




Trees

Part 8

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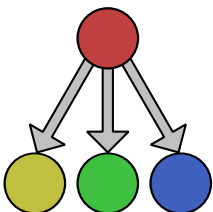
Introduction to Trees

Let the data grow

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Introduction to Trees

- In computer science, a tree is an abstract model of a hierarchical structure
- A tree consists of nodes with a parent-child relationship to zero *or more* nodes




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Some Applications

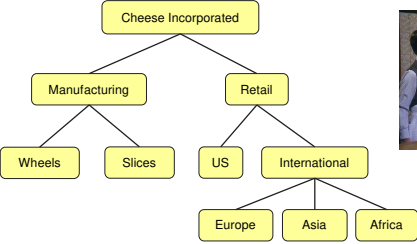
- Organizational charts
- Class hierarchy
- Disk directory and subdirectories
- Structure of a program



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Tree Example



```

graph TD
    CI[Cheese Incorporated] --> M[Manufacturing]
    CI --> R[Retail]
    M --> W[Wheels]
    M --> S[Slices]
    R --> US[US]
    R --> I[International]
    I --> E[Europe]
    I --> A[Asia]
    I --> AF[Africa]
  
```

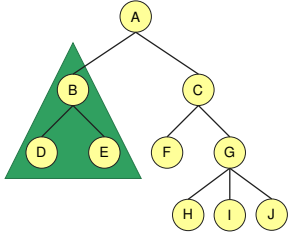


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Trees are Recursive

- Trees are recursive data structures
- They can be defined as smaller instances of trees
- So, using recursion is a natural approach



```

graph TD
    A((A)) --> B((B))
    A --> C((C))
    B --> D((D))
    B --> E((E))
    C --> F((F))
    C --> G((G))
    G --> H((H))
    G --> I((I))
    G --> J((J))
  
```

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Linked Lists vs. Trees

- **Linked Lists**
 - linear - accessing all elements is $O(n)$
 - nodes can only have one predecessor and/or one successor node
- **Trees**
 - nonlinear and hierarchical
 - nodes can have *multiple* successors but only one predecessor

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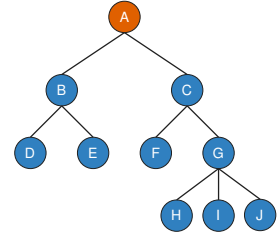
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Tree Terminology

- **Node**
 - just like in linked lists, the units of linked data are called nodes
 - usually contain data
- **Root**
 - starting point of the tree
 - no nodes link to it
 - e.g. A



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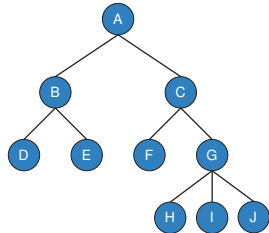
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Tree Terminology

- **Ancestor node**
 - predecessors
 - human-like lineage names: parent, grandparent, etc.
- **Descendant node**
 - successors
 - e.g. child, grandchild, great-grandchild, etc.



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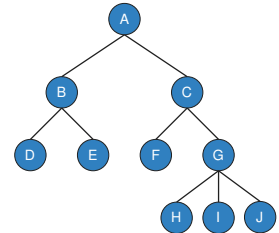
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Tree Terminology

- **Depth** of a node
 - # of ancestors to the root
 - e.g. depth of F is 2
- **Height** of a tree
 - maximum depth of any node
 - e.g. this tree is 3



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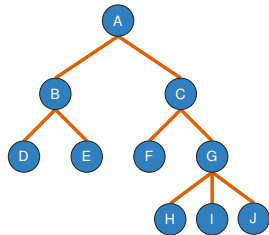
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Tree Terminology

- **Branch**
 - links between nodes
 - often unidirectional
- **Branching-factor**
 - max number of branches any node can have
 - can be 2 to more



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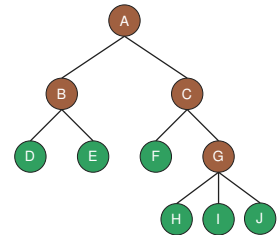
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Tree Terminology

- **Internal node**
 - node with at least one child
 - e.g. A, B, C, G
- **Leaf**
 - aka *external node*
 - node without children
 - e.g. D, E, F, H, I, J



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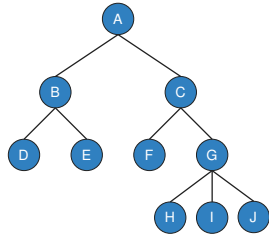
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Tree Terminology

- Size of the tree
 - total number of nodes
 - this tree has a size of 10



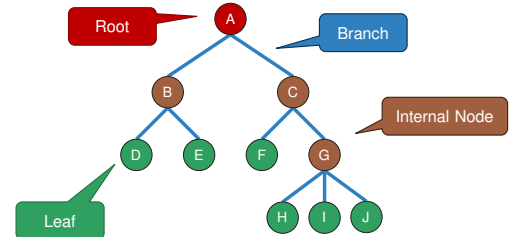
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Tree Terminology



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General Tree Node ADT

```
class Node
  public Object value; //Anything
  public Node[] branches;
end class
```

Array, or better, a linked list

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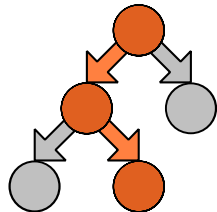


Climbing Down

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Tree Traversal

- A *tree traversal* visits the nodes of a tree in a systematic manner
- Given that trees can be defined into smaller and smaller subtrees, *recursion is an eloquent solution*



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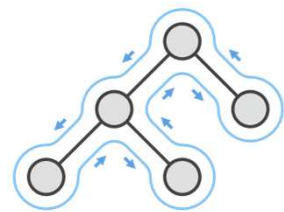
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Depth First Traversal

- If we continuously follow the tree to the left – this will result in *Euler Tour*
- We traverse the tree and pass through each node

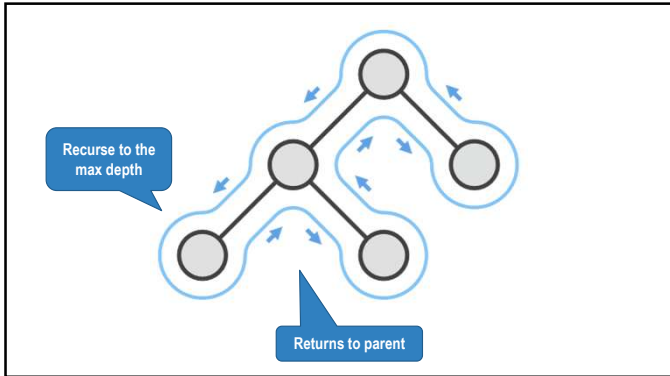


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Depth First Traversal

- Notice, in this case, that we tend to go the bottom first
- This is also known *depth-first traversal*

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Depth First Traversal

- This approach lends itself to recursion
- How?
 - root recurses into its children
 - each child recurses into each of its children

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Depth First Traversal

- A node is *visited* when its contents are analyzed
- Notice that we pass by each node going down and going up
- On either of these passes, we can visit the node

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Depth First Traversal

- This can be before or after its children are visited
- When the node is visited, when recursing the tree, has a **huge** impact on the algorithm

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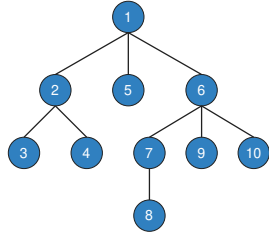
Depth-first: Preorder

- In a *preorder depth-first traversal*, a node is visited **before** its descendants
- In the image to the right, nodes will be visited in the order they are numbered

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Depth-first: Preorder

- Notice that each child was visited **after** its parent
- Some uses...
 - print a tree document
 - e.g. XML export



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Preorder Traversal Logic

```
function preorder
  this.visit()

  for each child c in this node
    c.preorder()
  end for
end function
```

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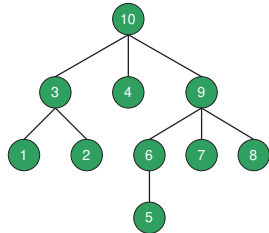
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Depth First: Postorder

- In a *postorder depth-first traversal*, a node is visited **after** its descendants
- Notice that each child was visited **before** its parent



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Depth First: Postorder

```
function postOrder
  for each child c in this node
    c.postOrder()
  end for

  this.visit()
end function
```

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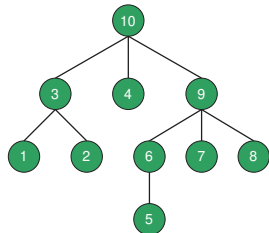
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Some Uses for Postorder

- Compute space used child nodes
- Calculate folder space
- Expression evaluation (an alternative to the stack algorithm)



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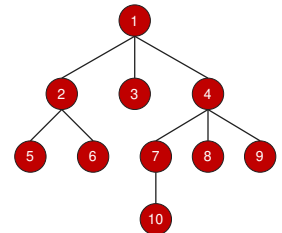
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Breadth-first Traversal

- In a *breadth-first traversal*, nodes are visited by their level in the tree
- So, the traversal, looks at all the nodes at depth 1, then at 2, etc...



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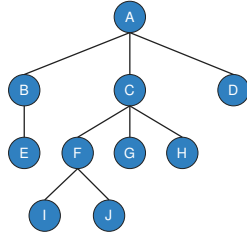
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Test Your Might

What is the order the nodes are visited using depth-first *pre-order* traversal?

ABECFIJGHD



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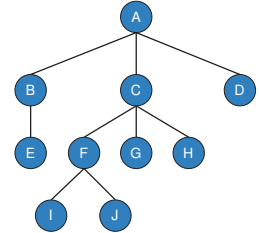
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Test Your Might

What is the order the nodes are visited using depth-first *post-order* traversal?

EBIJFGHCDA



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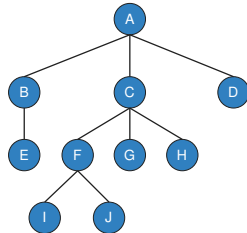
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Test Your Might

What is the order the nodes are visited using depth-first *breadth-first* traversal?

ABCDEFGHIJ



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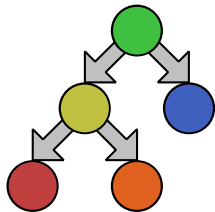
Binary Trees

The Power of Two!

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Binary Trees

- The most common tree used in data structures is in the style of the binary tree
- As the name implies, nodes in a binary tree only have two successors



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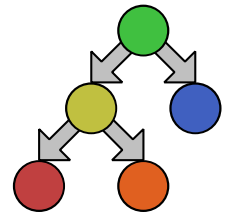
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Binary Trees

- We call the children of an internal node *left child* and *right child*



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Binary Trees

- Binary Trees are extremely useful in data structures
- The two branches allow for efficient branching and is ideal for binary operations
- Applications:
 - storing arithmetic expressions
 - decision processes
 - searching
 - sorting

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Binary Tree Node

```
class Node
  public Object value; //Can be anything
  public Node left;
  public Node right;
end class
```

Branches are much simpler

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Boolean Decision Tree



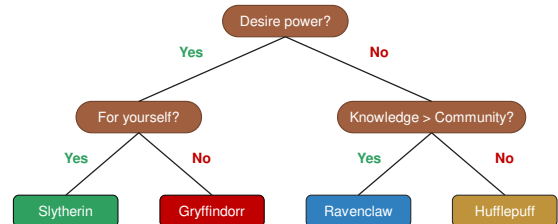
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Boolean Decision Tree



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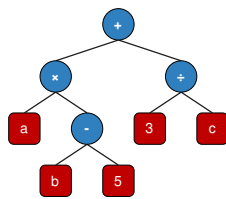
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Arithmetic Expression Tree

- Expressions can be represented with a tree
- How?
 - internal nodes: operators
 - leaves: operand



(a * (b - 5) + 3 / c)

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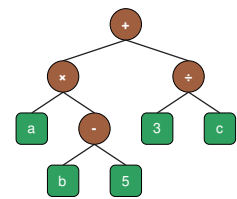
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Arithmetic Expression Tree

- It can be evaluated using a depth-first traversal
- ... notice that the node's children need a result before the node can be evaluated



(a * (b - 5) + 3 / c)

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Attributes of a Binary Tree

- $v = i + 1$
- $n = 2v - 1$
- $h \leq i$
- $h \leq (n - 1) / 2$
- $v \leq h$
- $h \geq \log_2 v$
- $h \geq \log_2 (n + 1) - 1$

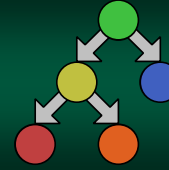
n	number of nodes
i	number of internal nodes
v	number of leaves
h	height of the tree

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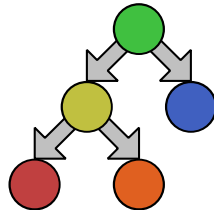
Depth-First Traversing Binary Trees

With simplicity, we have power!

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Depth-First Traversing

- Because of the simplicity of binary trees, we have a very useful structure for tree traversal
- We can only traverse left and right
- This gives **three** possibilities for a depth first search



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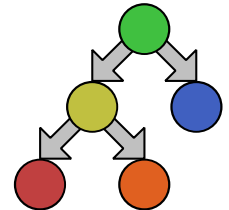
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Pre-order Depth-first Traversal

- When a *pre-order* depth-first traversal is performed, the node is visited **before** the right or left child
- Useful for copying a tree and printing trees



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Binary Pre-order Traversal Logic

```
function preOrder
  this.visit()

  if left isn't null then left.preOrder()
  if right isn't null then right.preOrder()
end function
```

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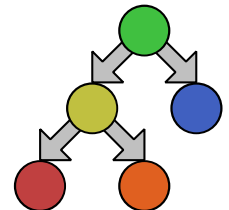
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In-order Depth-first Traversal

- In an *in-order* traversal a node is visited **after** its left branch and **before** its right branch
- In other words: recurse left, visit, then recurse right



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Binary In-order Traversal Logic

```
function inOrder
  if left isn't null then left.inOrder()

  this.visit()

  if right isn't null then right.inOrder()
end function
```

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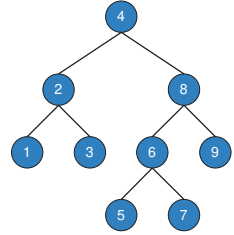
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Some In-order Applications

- Draw a binary tree
- Heap sorting
- Binary searching – $O(\log n)$ when sorted



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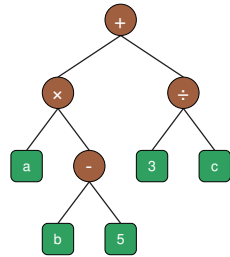
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In-order: Print Expressions

- In-order can be used to easily print an expression stored in a tree
- Print....
 - (then traverse left
 - the node's operator
 - traverse right then)



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In-order: Print Expressions

```
function print()
  if this is a leaf
    write this.value
  else
    write "("
    left.print()
    write this.operator ...can be stored in this.value
    right.print()
    write ")"
  end if
end function
```

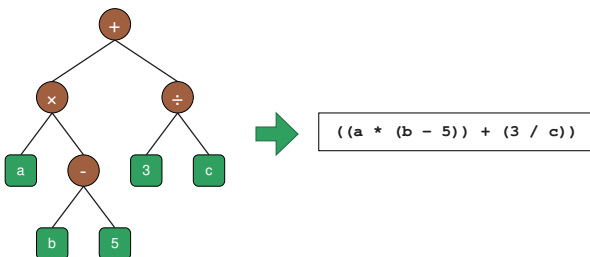
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In-order: Print Expressions



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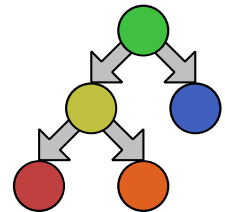
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Post-order Depth-first Traversal

- In a *post-order traversal* a node is evaluated after its left branch and after its right branch
- In other words: recurse left, recurse right, then visit



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Binary Post-order Traversal Logic

```
function postOrder
  if left isn't null then left.postOrder()
  if right isn't null then right.postOrder()

  this.visit()
end function
```

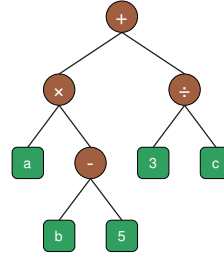
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Post-order: Evaluate Expressions



- A post-order traversal can be used to evaluate the tree
- Each recursive call (**left**, **right**) returns a value – the result of its calculation

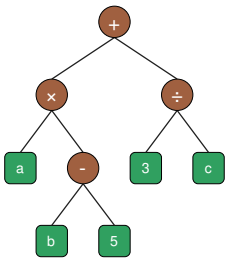
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Post-order: Evaluate Expressions



- The node then applies the operator to the two returned values (**left**, **right**)
- ... and then returns that value to its caller

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Post-order: Evaluate Expressions

```
function evaluate()
  if this is a leaf
    return this.value
  else
    x ← left.evaluate()
    y ← right.evaluate()
    ◊ ← this.operator ...can be stored in this.value
    return x ◊ y
  end if
end function
```

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