


Set Data Structures

Part 12

1



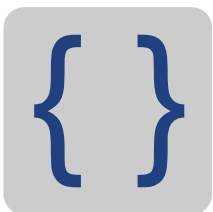
Importance of Sets

Organizing Information

2

Importance of Sets

- A *set* is an **unordered** collection of "objects"
- Sets are used in computer science in a wide variety of ways

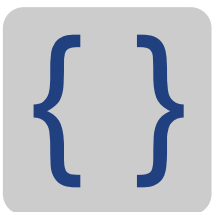


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3

Importance of Sets

- Depending on the attributes of the set, and how we use it, there are different approaches
- Programmers must choose the best model given how it will be used



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4

Set Review: Membership

- Set notation uses special symbols to denote if an object is a member of a set
- Below, the set V contains vegetables


potato $\in V$
bacon $\notin V$

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5

Set Review: Union

- A union of two sets combines all members of each set into a new one
- So, the result is two merged sets



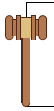
$A \cup B = \{ x \mid x \in A \text{ or } x \in B \}$

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6

Set Review: Intersection

- The intersection of two sets contains only those elements that are found in **both** sets
- So, the result is where the two sets overlap



$$A \cap B = \{ x \mid x \in A \text{ and } x \in B \}$$

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7

7

Set Review: Difference

- **Difference** (aka exclusion) removes all items found in set from another
- Typically, it is written either $A - B$ or $A \setminus B$



$$A - B = \{ x \mid x \in A \text{ and } x \notin B \}$$

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8

8

Set Review: Complement

- The **complement** of a set A , is all elements in the Universe, **not** in A
- Typically written as A' or A^c
- Alternatively written using an overbar



$$A' = \{ x \mid x \notin A \}$$

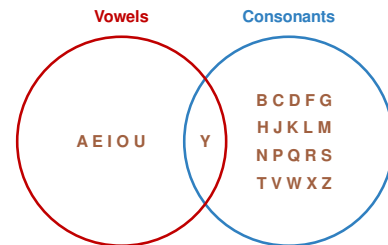
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9

9

Example: Letters in English

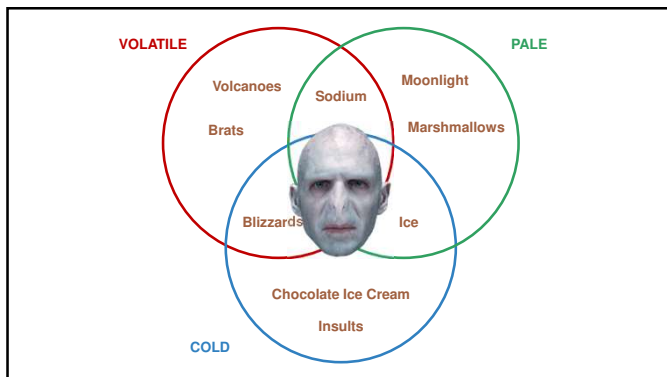


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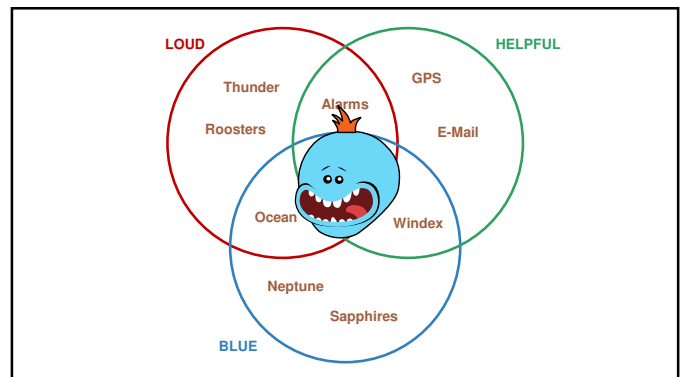
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10

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11



12

Multiset

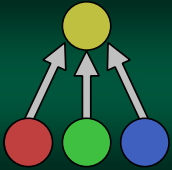


- A *multiset* is closely related to a set, but permits duplicate elements (as does a tuple)
- The Bag ADT, from the beginning of the semester, supports a basic multiset

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13

Union-Find



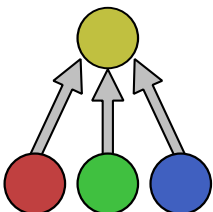
When it absolutely, positively has to be unioned overnight ... er... $O(1)$

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14

Union-Find

- Often there are situations where we have a number of known objects
- ... i.e. a finite countable set
- ... and we want to quickly union them together

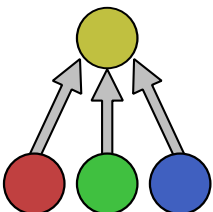


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15

Union-Find

- *Union-Find* data structure maintains a list of nodes *partitioned* into *disjoint* subsets
- e.g. $\{\{a\} \{b\} \{c\} \{d\} \{e\}\}$ is a full partition of $\{a,b,c,d,e\}$

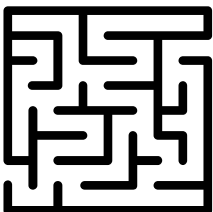


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16

Example Applications

- Kruskal's Tree-spanning Algorithm – uses a set of known edges
- Maze creation algorithms use it to track sets of connected paths – so the maze is always solvable

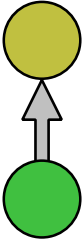


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17

The Approach

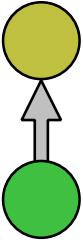
- Sets are stored as a variation of the classic tree
- However, in this approach **children link to their parents**
- So, the branches point **"backwards"** from standard trees – i.e. upwards



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18

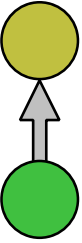
Arrangement of Nodes



- A parent node can have multiple children - this is **not** a binary tree!
- Every node - in the tree - is part of the same set

19

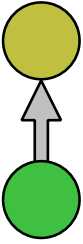
Arrangement of Nodes



- The root is called the *representative* of the set
- This node is **not** special nor is its value special
- ... it just happens to be the node the represents the set

20

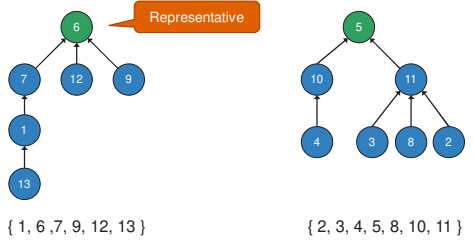
Arrangement of Nodes



- Any node can "find" its representative – by following the upwards branches
- If any two nodes have the same representative, then they are in the same set

21

Examples



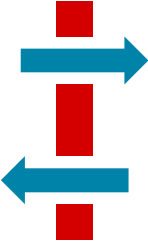
{ 1, 6, 7, 9, 12, 13 }

{ 2, 3, 4, 5, 8, 10, 11 }

22

Union-Find Operations

- The Union-Find data structure contains 3 operations that handle sets
- All 3 can modify the structure of the tree – which automatically optimizes itself into $O(1)$



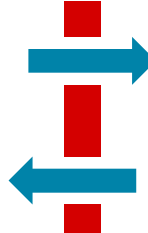
23

Union-Find Interface

```
public class UnionFind
{
    void makeSet(object value)
    void union(object a, object b)
    object find(object a)
}
```

24

Makeset




- Creates a set (within the Union-Find instance) that contains only a single element
- Also known as a *singleton*
- Essentially, this is the same as an "add" for the ADT

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25

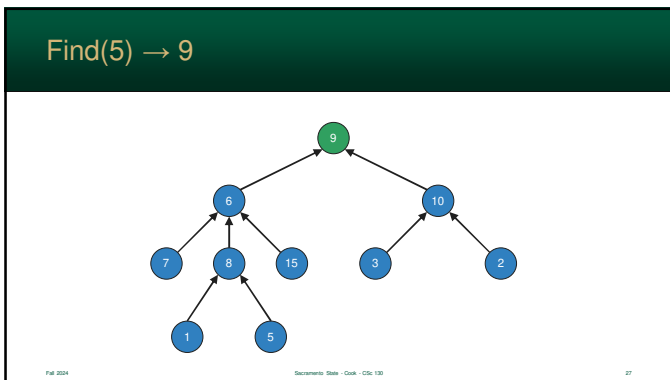
Find



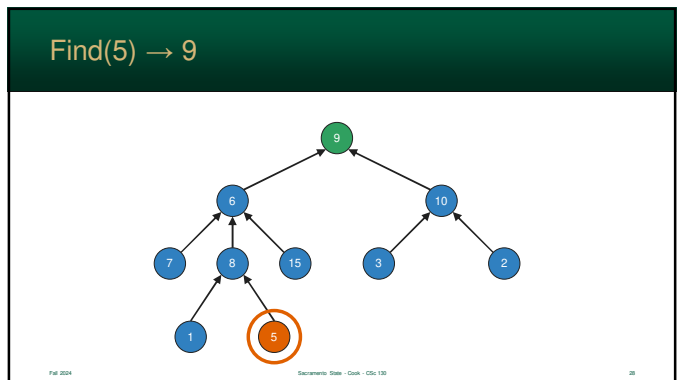
- The *find operation* starts with a node and locates its *representative*
- Travels up the tree until it finds a node without a link (which is the representative)

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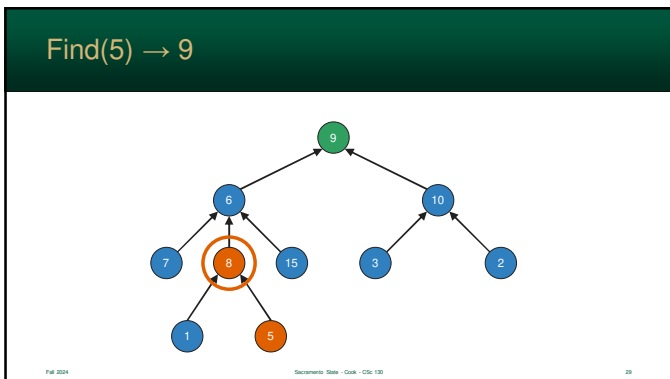
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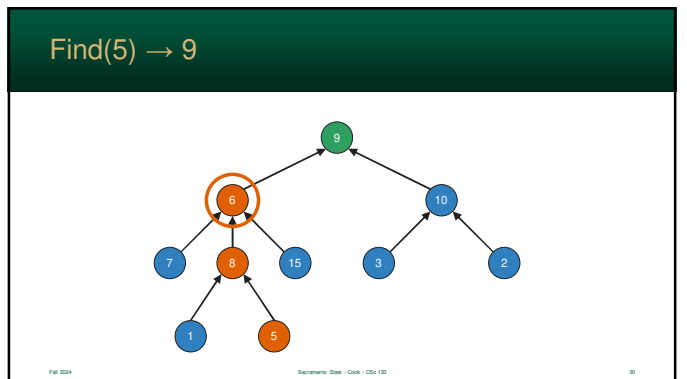
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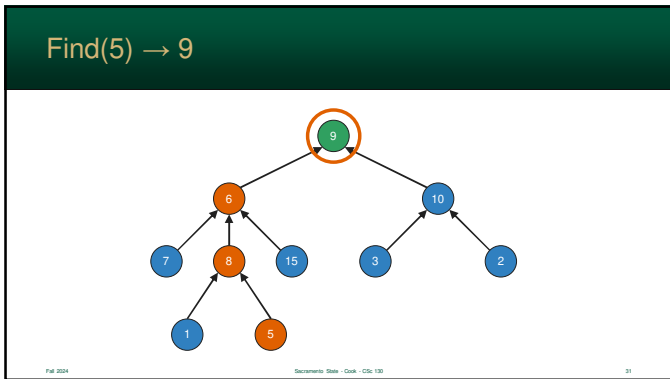
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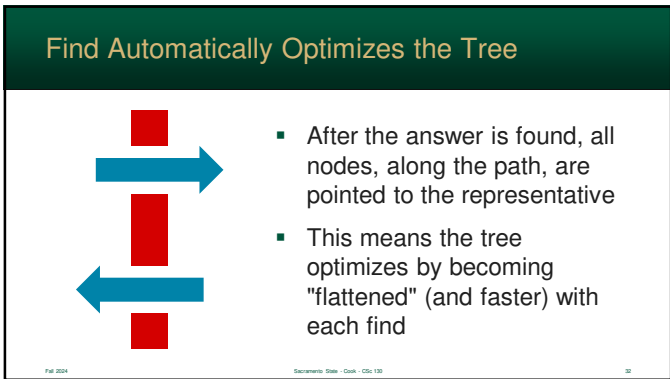
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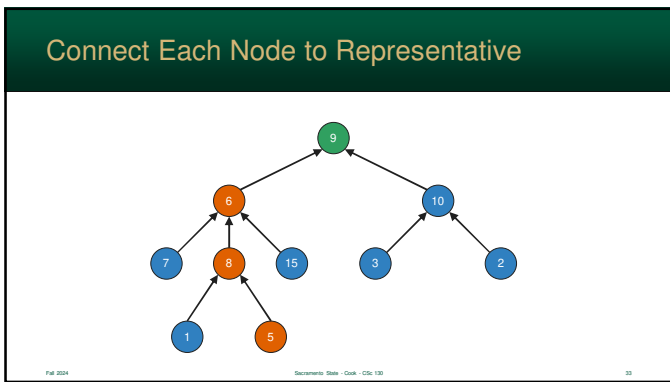
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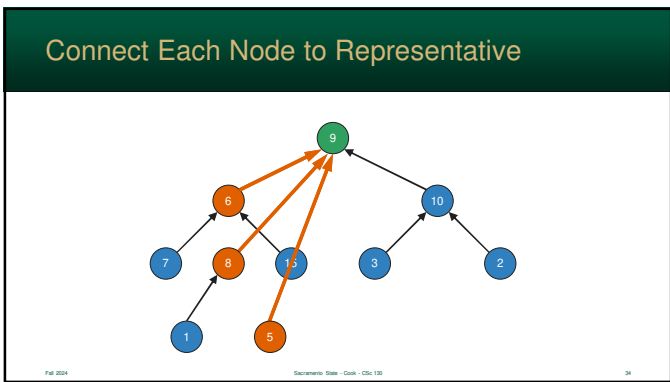
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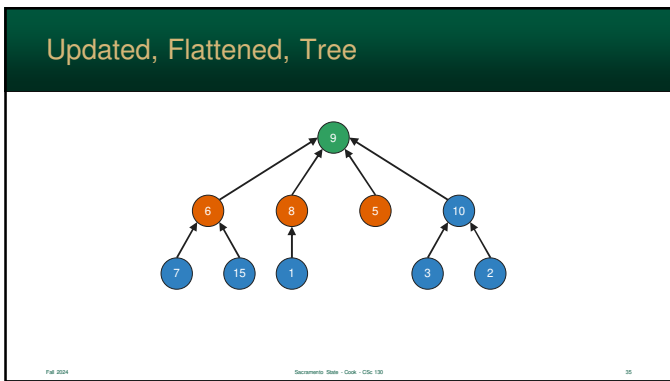
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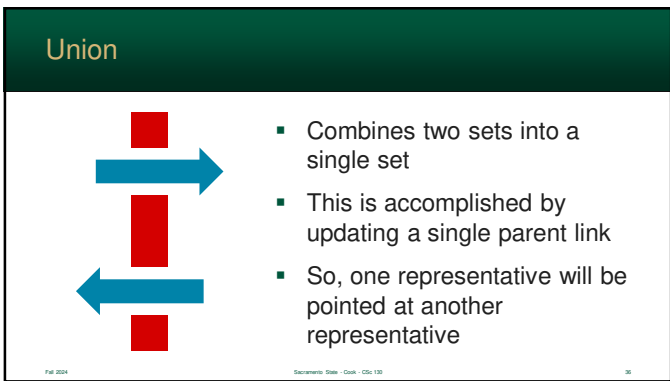
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34

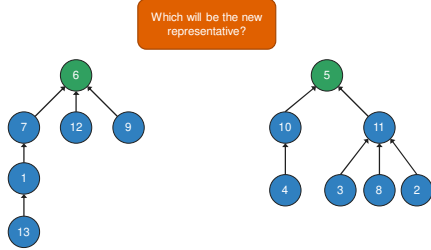


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36

So, which one do we choose?



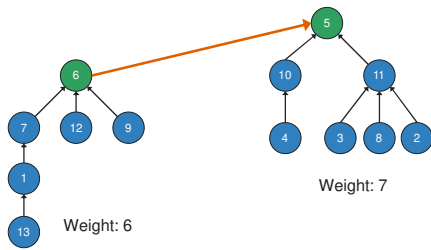
37

Approach #1: Union by Weight

- Using *Union by Weight*, the tree with the fewer nodes (not edge weight) is made a subtree of the larger tree
- This helps create a more balanced union in terms of weight
- Again, the Find operation will optimize the tree

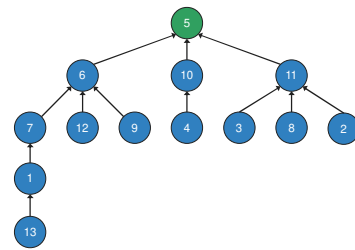
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So, which one do we choose?



39

Approach #1: Result



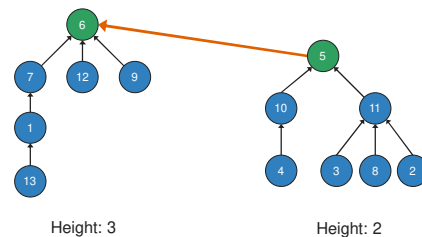
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Approach #2: Union By Height

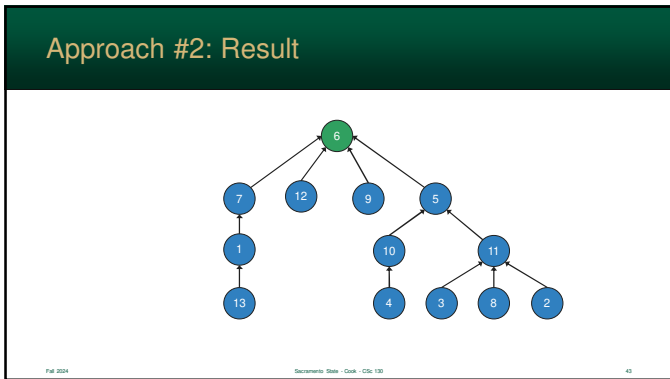
- Using *Union by Height*, the tree with the smaller height is made a subtree of the taller tree
- This helps create a more balanced union in terms of height
- But, Find, will fix this automatically

41

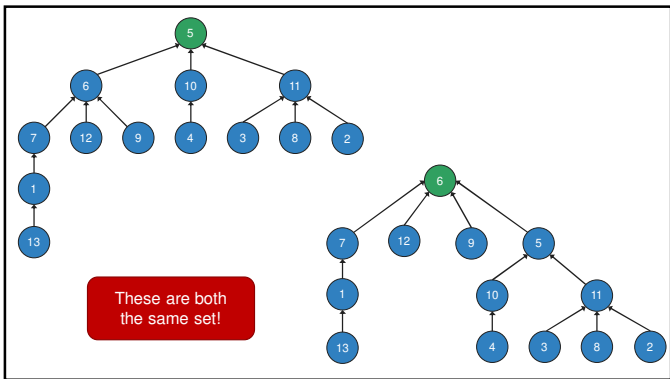
Approach #2: Union By Height



42



43



44

Bit Vectors

Sets and Bits

45

Bit Vectors

- A *bit vector* can store **finite**, **countable** sets using bits
- Also known as a *bit array*, *bit set*, and *bit map*
- Compact format that can perform a set operations with a single operation (**fast!**)

46

Bit Vectors

- Each object in the universe is represented as a **single** bit in the string of bits
- If $x \in A$, then the bit is **1**, otherwise **0**

47

Example 1

- $U = \{ \text{fry, leela, zoidberg, bender, hermes} \}$
- $A = \{ \text{fry, leela, bender} \}$

U = 11111
A = 11010

48

Example 1

- $U = \{ \text{rick, beth, jerry, morty, summer} \}$
- $A = \{ \text{rick, morty} \}$

$U = 11111$
 $A = 10010$

49

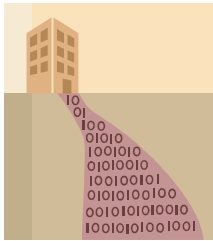
Example 2

- $U = \{ 2, 3, 5, 7, 11, 13, 17, 19 \}$
- $A = \{ 3, 5, 11, 19 \}$

$U = 11111111$
 $A = 01101001$

50

Why this is useful



- Computers can easily perform **and** & **or** operations on bytes (or multiple bytes)
- This means set operations can be performed amazingly fast

51

Let's look at the definitions again...

- The definitions of union and intersection are nearly identical
- The relationship between the elements is defined using an **AND** or **OR**

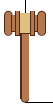


$$A \cup B = \{ x \mid x \in A \text{ or } x \in B \}$$
$$A \cap B = \{ x \mid x \in A \text{ and } x \in B \}$$

52

Let's look at the definitions again...

- We can apply **bit-wise-and** & **bit-wise-or**
- It will apply the operation to each of the bits in matching columns



$$A \cup B = \{ x \mid x \in A \text{ or } x \in B \}$$
$$A \cap B = \{ x \mid x \in A \text{ and } x \in B \}$$

53

Let's look at the definitions again...

- So, each bit in **A** will be compared to its matching bit in **B**
- Bit match can do sets!



$$A \cup B = \{ x \mid x \in A \text{ or } x \in B \}$$
$$A \cap B = \{ x \mid x \in A \text{ and } x \in B \}$$

54

Example: Union (using or)

$U = \{a, b, c, d, e, f, g\}$
 $A = \{b, c, d\} = 0111000$
 $B = \{d, e, f\} = 0001110$

0111000
or 0001110
 $0111110 = \{b, c, d, e, f\}$

55

Example: Intersection (using and)

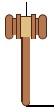
$U = \{a, b, c, d, e, f, g\}$
 $A = \{b, c, d\} = 0111000$
 $B = \{d, e, f\} = 0001110$

0111000
and 0001110
 $0001000 = \{d\}$

56

Complement

- Then, how do we do a complement of a set A ?
- We must flip all the bits from 1 to 0, and 0 to 1
- We can use a *binary-not* or the *XOR* operation



$$A' = \{x \mid x \notin A\}$$

57

Example: Complement (using not)

$U = \{a, b, c, d, e, f, g\}$
 $A = \{b, c, d\} = 0111000$

not 0111000
 $1000111 = \{a, e, f, g\}$

58

Exclusion

- Finally, how do we do set difference?
- The "subtract" operator will not work
- Let's look at the definition a bit more closely



$$A - B = \{x \mid x \in A \text{ and } x \notin B\}$$

59

Exclusion

- It's essentially the definition of *intersection*
- Except, the second operand is the definition of *complement*.



$$A - B = \{x \mid x \in A \text{ and } x \notin B\}$$

60

Example: Intersection (using and)

```
U = {a, b, c, d, e, f, g}
A = {b, c, d, f} = 0111010
B = {d, e, f} = 0001110
B' = not 0001110 = 1110001

  0111010
and 1110001
  0110000 = {b, c}
```

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61

61

Java/C Code

```
Intersection : a & b
Union        : a | b
Complement  : ~a
Exclusion     : a & ~b
```

&& and || are Boolean operators. These are bit-wise.

The tilde ~ is a bitwise not.

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62

62